        A Session Initiation Protocol (SIP) Load-Control Event Package

Abstract

   This specification defines a load-control event package for the
   Session Initiation Protocol (SIP). It allows SIP entities to
   distribute load-filtering policies to other SIP entities in the
   network. The load-filtering policies contain rules to throttle calls
   from a specific user or based on their source or destination domain,
   telephone number prefix. The mechanism helps to prevent signaling
   overload and complements feedback-based SIP overload control efforts.

Status of This Memo

   This is an Internet Standards Track document.

   This document is a product of the Internet Engineering Task Force
   (IETF). It represents the consensus of the IETF community. It has
   received public review and has been approved for publication by the
   Internet Engineering Steering Group (IESG). Further information on
   Internet Standards is available in Section 2 of RFC 5741.

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   and how to provide feedback on it may be obtained at
   http://www.rfc-editor.org/info/rfc7200.

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1. Introduction

   SIP load-control mechanisms are needed to prevent congestion collapse
   [RFC6357] in cases of SIP server overload [RFC5390]. There are two
   types of load-control approaches. In the first approach, feedback
   control, SIP servers provide load limits to upstream servers, to
   reduce the incoming rate of all SIP requests [SIP-OVERLOAD]. These
   upstream servers then drop or delay incoming SIP requests. Feedback
   control is reactive and affects signaling messages that have already
   been issued by user agent clients. This approach works well when SIP
   proxy servers in the core networks (core proxy servers) or
   destination-specific SIP proxy servers in the edge networks (edge
   proxy servers) are overloaded. By their nature, they need to
   distribute rate, drop, or window information to all upstream SIP
   proxy servers and normally affect all calls equally, regardless of
   destination.

   This specification proposes an additional, complementary load-control
   mechanism, called "load filtering". It is most applicable for
   situations where a traffic surge and its source/destination
   distribution can be predicted in advance. In those cases, network
   operators create load-filtering policies that indicate calls to
   specific destinations or from specific sources should be rate-limited
   or randomly dropped. These load-filtering policies are then
   distributed to SIP servers and possibly SIP user agents that are
   likely to generate calls to the affected destinations or from the
   affected sources. Load filtering works best if it prevents calls as
   close to the originating user agent clients as possible. The
   applicability of SIP load filtering can also be extended beyond
   overload control, e.g., to implement service level agreement
   commitments.

2. Conventions

   The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT",
   "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this
   document are to be interpreted as described in [RFC2119].

3. SIP Load-Filtering Overview

3.1. Load-Filtering Policy Format

   Load-filtering policies are specified by sets of rules. Each rule
   contains both load-filtering conditions and actions. The load-
   filtering conditions define identities of the targets to be filtered
   (Section 5.3.1). For example, there are two typical resource limits
   in a possible overload situation, i.e., human destination limits
   (number of call takers) and node capacity limits. The load-filtering
   targets in these two cases can be the specific callee numbers or the
   destination domain corresponding to the overload. Load-filtering
   conditions also indicate the specific message type to be matched
   (Section 5.3.2), with which target SIP entity the filtering policy is
   associated (Section 5.3.3), and the period of time when the filtering
   policy should be activated and deactivated (Section 5.3.4). Load-
   filtering actions describe the desired control functions such as
   keeping the request rate below a specified level (Section 5.4).

3.2. Load-Filtering Policy Computation

   When computing the load-filtering policies, one needs to take into
   consideration information such as overload time, scope and network
   topology, as well as service policies. It is also important to make
   sure that there is no resource allocation loop and that server
   capacity is allocated in a way that both prevents overload and
   maximizes effective throughput (commonly called goodput). In some
   cases, in order to better utilize system resources, it may be
   preferable to employ an algorithm that dynamically computes the load-
   filtering policies based on currently observed server load status,
   rather than using a purely static filtering policy assignment. The
   computation algorithm for load-filtering policies is beyond the scope
   of this specification.

3.3. Load-Filtering Policy Distribution

   For distributing load-filtering policies, this specification defines
   the SIP event package for load control, which is an "instantiation"
   of the generic SIP event notification framework [RFC6665]. This
   specification also defines the XML schema of a load-control document
   (Section 5), which is used to encode load-filtering policies.

   In order for load-filtering policies to be properly distributed, each
   capable SIP entity in the network subscribes to the SIP load-control
   event package of each SIP entity to which it sends signaling
   requests. A SIP entity that accepts subscription requests is called
   a "notifier" (Section 4.6). Subscription is initiated and maintained
   during normal server operation. The subscription of neighboring SIP

entities needs to be persistent, as described in Sections 4.1 and 4.2
of [RFC6665]. The refresh procedure is described in Section 4.7
below. Subscribers may terminate the subscription if they have not
received notifications for an extended time period, and can
resubscribe if they determine that signaling with the notifier
becomes active again.

An example architecture is shown in Figure 1 to illustrate SIP load-
filtering policy distribution. This scenario consists of two
networks belonging to Service Provider A and Service Provider B,
respectively. Each provider's network is made up of two SIP core
proxy servers and four SIP edge proxy servers. The core proxy
servers and edge proxy servers of Service Provider A are denoted as
CPa1 to CPa2 and EPa1 to EPa4; the core proxy servers and edge proxy
servers of Service Provider B are denoted as CPb1 to CPb2 and EPb1 to
EPb4.

```
+----------+   +----------+   +----------+   +----------+
|          |   |          |   |          |   |          |
|   EPa1   |   |   EPa2   |   |   EPa3   |   |   EPa4   |
|          |   |          |   |          |   |          |
+----------+   +----------+   +----------+   +----------+
       \         /                 \         /
        \       /                   \       /
         \     /                     \     /
      +----------+               +----------+
      |          |               |          |
      |   CPa1   |---------------|   CPa2   |
      |          |               |          |
      +----------+               +----------+
           |                          |
 Service   |                          |
 Provider A|                          |
           |                          |
 ==========|==========================|=========================
           |                          |
 Service   |                          |
 Provider B|                          |
           |                          |
      +----------+               +----------+
      |          |               |          |
      |   CPb1   |---------------|   CPb2   |
      |          |               |          |
      +----------+               +----------+
       /     \                     /     \
      /       \                   /       \
     /         \                 /         \
+----------+   +----------+   +----------+   +----------+
|          |   |          |   |          |   |          |
|   EPb1   |   |   EPb2   |   |   EPb3   |   |   EPb4   |
|          |   |          |   |          |   |          |
+----------+   +----------+   +----------+   +----------+
```
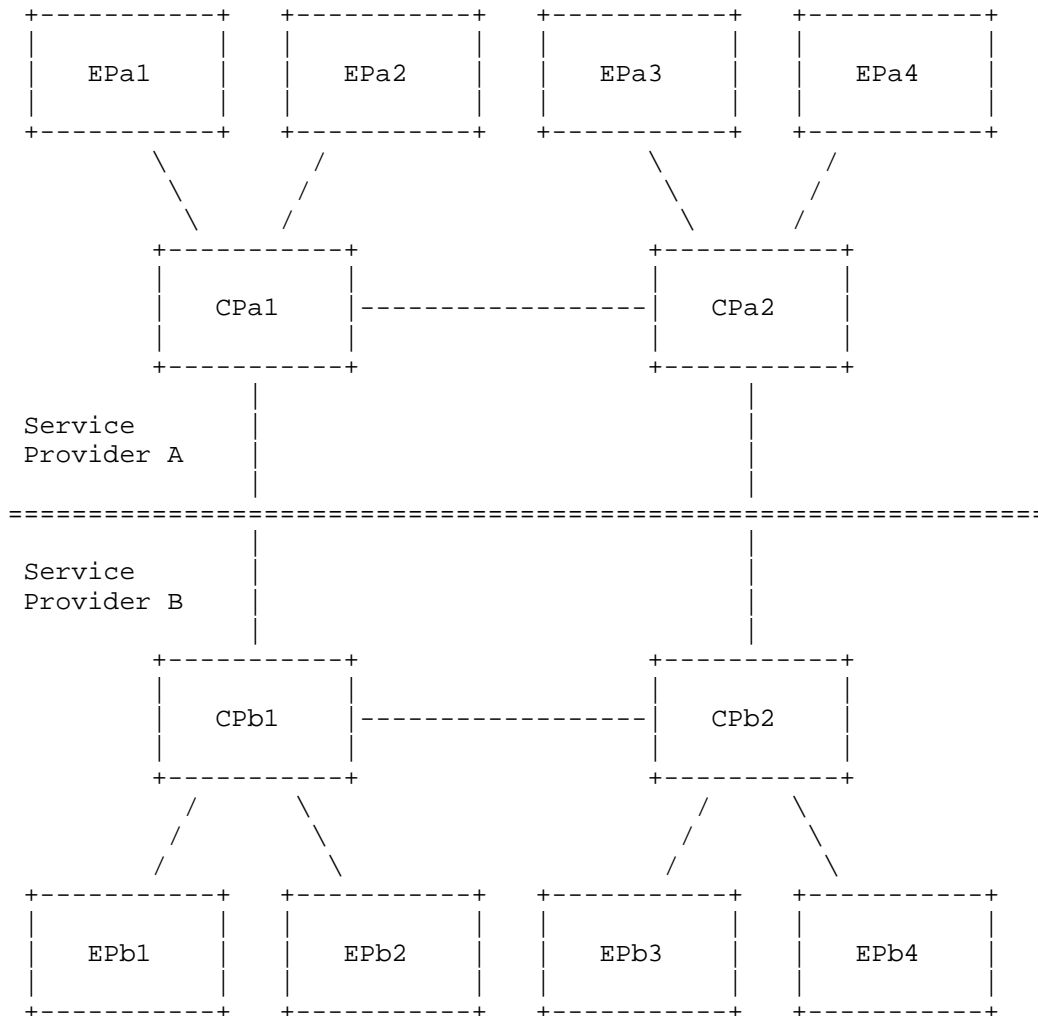
            Figure 1: Example Network Scenario Using SIP Load-Control Event
                              Package Mechanism

   During the initialization stage, the proxy servers first identify all
   their outgoing signaling neighbors and subscribe to them. Service
   providers can provision neighbors, or the proxy servers can
   incrementally learn who their neighbors are by inspecting signaling
   messages that they send and receive. Assuming all signaling
   relationships in Figure 1 are bidirectional, after this
   initialization stage, each proxy server will be subscribed to all its
   neighbors.

Case I: EPa1 serves a TV program hotline and decides to limit the
total number of incoming calls to the hotline to prevent an overload.
To do so, EPa1 sends a notification to CPa1 with the specific hotline
number, time of activation, and total acceptable call rate.
Depending on the load-filtering policy computation algorithm, CPa1
may allocate the received total acceptable call rate among its
neighbors, namely, EPa2, CPa2, and CPb1, and notify them about the
resulting allocation along with the hotline number and the activation
time. CPa2 and CPb1 may perform further allocation among their own
neighbors and notify the corresponding proxy servers. This process
continues until all edge proxy servers in the network have been
informed about the event and have proper load-filtering policies
configured.

In the above case, the network entity where load-filtering policy is
first introduced is the SIP server providing access to the resource
that creates the overload situation. In other cases, the network
entry point of introducing load-filtering policy could also be an
entity that hosts this resource. For example, an operator may host
an application server that performs toll-free-number ("800 number")
translation services. The application server itself may be a SIP
proxy server or a SIP Back-to-Back User Agent (B2BUA). If one of the
toll-free numbers hosted at the application server creates the
overload condition, the load-filtering policies can be introduced
from the application server and then propagated to other SIP proxy
servers in the network.

Case II: A hurricane affects the region covered by CPb2, EPb3, and
EPb4. All three of these SIP proxy servers are overloaded. The
rescue team determines that outbound calls are more valuable than
inbound calls in this specific situation. Therefore, EPb3 and EPb4
are configured with load-filtering policies to accept more outbound
calls than inbound calls. CPb2 may be configured the same way or
receive dynamically computed load-filtering policies from EPb3 and
EPb4. Depending on the load-filtering policy computation algorithm,
CPb2 may also send out notifications to its outside neighbors, namely
CPb1 and CPa2, specifying a limit on the acceptable rate of inbound
calls to CPb2's responsible domain. CPb1 and CPa2 may subsequently
notify their neighbors about limiting the calls to CPb2's area. The
same process could continue until all edge proxy servers are notified
and have load-filtering policies configured.

Note that this specification does not define the provisioning
interface between the party who determines the load-filtering policy
and the network entry point where the policy is introduced. One of
the options for the provisioning interface is the Extensible Markup
Language (XML) Configuration Access Protocol (XCAP) [RFC4825].

3.4. Applicable Network Domains

   This specification MUST be applied inside a "Trust Domain". The
   concept of a Trust Domain is similar to that defined in [RFC3324] and
   [RFC3325]. A Trust Domain, for the purpose of SIP load filtering, is
   a set of SIP entities such as SIP proxy servers that are trusted to
   exchange load-filtering policies defined in this specification. In
   the simplest case, a Trust Domain is a network of SIP entities
   belonging to a single service provider who deploys it and accurately
   knows the behavior of those SIP entities. Such simple Trust Domains
   may be joined to form larger Trust Domains by bilateral agreements
   between the service providers of the SIP entities.

   The key requirement of a Trust Domain for the purpose of SIP load
   filtering is that the behavior of all SIP entities within a given
   Trust Domain is known to comply to the following set of
   specifications.

   o SIP entities in the Trust Domain agree on the mechanisms used to
      secure the communication among SIP entities within the Trust
      Domain.

   o SIP entities in the Trust Domain agree on the manner used to
      determine which SIP entities are part of the Trust Domain.

   o SIP entities in the Trust Domain are compliant to SIP [RFC3261].

   o SIP entities in the Trust Domain are compliant to SIP-Specific
      Event Notification[RFC6665].

   o SIP entities in the Trust Domain are compliant to this
      specification.

   o SIP entities in the Trust Domain agree on what types of calls can
      be affected by this SIP load-filtering mechanism. For example,
      <call-identity> condition elements (Section 5.3.1) <one> and
      <many> might be limited to describe within certain prefixes.

   o SIP entities in the Trust Domain agree on the destinations to
      which calls may be redirected when the "redirect" action
      (Section 5.4) is used. For example, the URI might have to match a
      given set of domains.

   SIP load filtering is only effective if all neighbors that are
   possible signaling sources participate and enforce the designated
   load-filtering policies. Otherwise, a single non-conforming neighbor
   could make all filtering efforts useless by pumping in excessive
   traffic to overload the server. Therefore, the SIP server that

distributes load-filtering policies needs to take countermeasures
towards any non-conforming neighbors. A simple method is to reject
excessive requests with 503 "Service Unavailable" response messages
as if they were obeying the rate. Considering the rejection costs, a
more complicated but fairer method would be to allocate at the
overloaded server the same amount of processing to the combination of
both normal processing and rejection as the overloaded server would
devote to processing requests for a conforming upstream SIP server.
These approaches work as long as the total rejection cost does not
overwhelm the entire server resources. In addition, SIP servers need
to handle message prioritization properly while performing load
filtering, which is described in Section 4.8.

4. Load-Control Event Package

   The SIP load-filtering mechanism defines a load-control event package
   for SIP based on [RFC6665].

4.1. Event Package Name

   The name of this event package is "load-control". This name is
   carried in the Event and Allow-Events header, as specified in
   [RFC6665].

4.2. Event Package Parameters

   No package-specific event header field parameters are defined for
   this event package.

4.3. SUBSCRIBE Bodies

   This specification does not define the content of SUBSCRIBE bodies.
   Future specifications could define bodies for SUBSCRIBE messages, for
   example, to request specific types of load-control event
   notifications.

   A SUBSCRIBE request sent without a body implies the default
   subscription behavior as specified in Section 4.7.

4.4. SUBSCRIBE Duration

   The default expiration time for a subscription to load-filtering
   policy is one hour. Since the desired expiration time may vary
   significantly for subscriptions among SIP entities with different
   signaling relationships, the subscribers and notifiers are
   RECOMMENDED to explicitly negotiate appropriate subscription duration
   when knowledge about the mutual signaling relationship is available.

4.5. NOTIFY Bodies

   The body of a NOTIFY request in this event package contains load-
   filtering policies. The format of the NOTIFY request body MUST be in
   one of the formats defined in the Accept header field of the
   SUBSCRIBE request or be the default format, as specified in
   [RFC6665]. The default data format for the NOTIFY request body of
   this event package is "application/load-control+xml" (defined in
   Section 5). This means that when a NOTIFY request body exists but no
   Accept header field is specified in a SUBSCRIBE request, the NOTIFY
   request body MUST contain content conforming to the "application/
   load-control+xml" format.

4.6. Notifier Processing of SUBSCRIBE Requests

   The notifier accepts a new subscription or updates an existing
   subscription upon receiving a valid SUBSCRIBE request.

   If the identity of the subscriber sending the SUBSCRIBE request is
   not allowed to receive load-filtering policies, the notifier MUST
   return a 403 "Forbidden" response.

   If none of the media types specified in the Accept header of the
   SUBSCRIBE request are supported, the notifier SHOULD return a 406
   "Not Acceptable" response.

4.7. Notifier Generation of NOTIFY Requests

   A notifier MUST send a NOTIFY request with its current load-filtering
   policy to the subscriber upon successfully accepting or refreshing a
   subscription. If no load-filtering policy needs to be distributed
   when the subscription is received, the notifier SHOULD sent a NOTIFY
   request without a body to the subscriber. The content-type header
   field of this NOTIFY request MUST indicate the correct body format as
   if the body were present (e.g., "application/load-control+xml").
   Notifiers are likely to send NOTIFY requests without a body when a
   subscription is initiated for the first time, e.g., when a SIP entity
   is just introduced, because there may be no planned events that
   require load filtering at that time. A notifier SHOULD generate
   NOTIFY requests each time the load-filtering policy changes, with the
   maximum notification rate not exceeding values defined in
   Section 4.10.

4.8. Subscriber Processing of NOTIFY Requests

   The subscriber is the load-filtering server that enforces load-
   filtering policies received from the notifier. The way subscribers
   process NOTIFY requests depends on the load-filtering policies

conveyed in the notifications. Typically, load-filtering policies
consist of rules specifying actions to be applied to requests
matching certain conditions. A subscriber receiving a notification
first installs these rules and then enforces corresponding actions on
requests matching those conditions, for example, limiting the sending
rate of call requests destined for a specific callee.

In the case when load-filtering policies specify a future validity,
it is possible that when the validity time arrives, the subscription
to the specific notifier that conveyed the rules has expired. In
this case, it is RECOMMENDED that the subscriber re-activate its
subscription with the corresponding notifier. Regardless of whether
or not this re-activation of subscription is successful, when the
validity time is reached, the subscriber SHOULD enforce the
corresponding rules.

Upon receipt of a NOTIFY request with a Subscription-State header
field containing the value "terminated", the subscription status with
the particular notifier will be terminated. Meanwhile, subscribers
MUST also terminate previously received load-filtering policies from
that notifier.

The subscriber MUST discard unknown bodies. If the NOTIFY request
contains several bodies, none of them being supported, it SHOULD
unsubscribe unless it has knowledge that it will possibly receive
NOTIFY requests with supported bodies from that notifier. A NOTIFY
request without a body indicates that no load-filtering policies need
to be updated.

When the subscriber enforces load-filtering policies, it needs to
prioritize requests and select those requests that need to be
rejected or redirected. This selection is largely a matter of local
policy. It is expected that the subscriber will follow local policy
as long as the result in reduction of traffic is consistent with the
overload algorithm in effect at that node. Accordingly, the
normative behavior described in the next three paragraphs should be
interpreted with the understanding that the subscriber will aim to
preserve local policy to the fullest extent possible.

o The subscriber SHOULD honor the local policy for prioritizing SIP
   requests such as policies based on message type, e.g., INVITEs
   versus requests associated with existing sessions.

o The subscriber SHOULD honor the local policy for prioritizing SIP
   requests based on the content of the Resource-Priority header
   (RPH, [RFC4412]). Specific (namespace.value) RPH contents may
   indicate high-priority requests that should be preserved as much

      as possible during overload. The RPH contents can also indicate a
      low-priority request that is eligible to be dropped during times
      of overload.

   o The subscriber SHOULD honor the local policy for prioritizing SIP
      requests relating to emergency calls as identified by the sos URN
      [RFC5031] indicating an emergency request.

   A local policy can be expected to combine both the SIP request type
   and the prioritization markings and SHOULD be honored when overload
   conditions prevail.

4.9. Handling of Forked Requests

   Forking is not applicable when this load-control event package
   mechanism is used within a single-hop distance between neighboring
   SIP entities. If communication scope of the load-control event
   package mechanism is among multiple hops, forking is also not
   expected to happen because the subscription request is addressed to a
   clearly defined SIP entity. However, in the unlikely case when
   forking does happen, the load-control event package only allows the
   first potential dialog-establishing message to create a dialog, as
   specified in Section 5.4.9 of [RFC6665].

4.10. Rate of Notifications

   The rate of notifications is unlikely to be of concern for this local
   control event package mechanism when it is used in a non-real-time
   mode for relatively static load-filtering policies. Nevertheless, if
   a situation does arise in which a rather frequently used load
   filtering policy update is needed, it is RECOMMENDED that the
   notifier not generate notifications at a rate higher than once per
   second in all cases, in order to avoid the NOTIFY request itself
   overloading the system.

4.11. State Delta

   It is likely that updates to specific load-filtering policies are
   made by changing only part of the policy parameters (e.g., acceptable
   request rate or percentage, but not matching identities). This will
   typically be because the utilization of a resource subject to
   overload depends upon dynamic unknowns such as holding time and the
   relative distribution of offered loads over subscribing SIP entities.
   The updates could originate manually or be determined automatically
   by an algorithm that dynamically computes the load-filtering policies
   (Section 3.2). Another factor that is usually not known precisely or

needs to be computed automatically is the duration of the event
requiring load filtering. Therefore, it would also be common for the
validity to change frequently.

This event package allows the use of state delta as in [RFC6665] to
accommodate frequent updates of partial policy parameters. For each
NOTIFY transaction in a subscription, a version number that increases
by exactly one MUST be included in the NOTIFY request body when the
body is present. When the subscriber receives a state delta, it
associates the partial updates to the particular policy by matching
the appropriate rule id (Appendix D). If the subscriber receives a
NOTIFY request with a version number that is increased by more than
one, it knows that it has missed a state delta and needs to ask for a
full state snapshot. Therefore, the subscriber ignores that NOTIFY
request containing the state delta, and resends a SUBSCRIBE request
to force a NOTIFY request containing a complete state snapshot.

5. Load-Control Document

5.1. Format

A load-control document is an XML document that describes the load-
filtering policies. It inherits and enhances the common policy
document defined in [RFC4745]. A common policy document contains a
set of rules. Each rule consists of three parts: conditions,
actions, and transformations. The conditions part is a set of
expressions containing attributes such as identity, domain, and
validity time information. Each expression evaluates to TRUE or
FALSE. Conditions are matched on "equality" or "greater than" style
comparison. There is no regular expression matching. Conditions are
evaluated on receipt of an initial SIP request for a dialog or
standalone transaction. If a request matches all conditions in a
rule set, the action part and the transformation part are consulted
to determine the "permission" on how to handle the request. Each
action or transformation specifies a positive grant to the policy
server to perform the resulting actions. Well-defined mechanism are
available for combining actions and transformations obtained from
more than one sources.

5.2. Namespace

The namespace URI for elements defined by this specification is a
Uniform Resource Namespace (URN) ([RFC2141]), using the namespace
identifier "ietf" defined by [RFC2648] and extended by [RFC3688].
The URN is as follows:

urn:ietf:params:xml:ns:load-control

5.3. Conditions

   [RFC4745] defines three condition elements: <identity>, <sphere>, and
   <validity>. This specification defines new condition elements and
   reuses the <validity> element. The <sphere> element is not used.

5.3.1. Call Identity

   Since the problem space of this specification is different from that
   of [RFC4745], the [RFC4745] <identity> element is not sufficient for
   use with load filtering. First, load filtering may be applied to
   different identities contained in a request, including identities of
   both the receiving entity and the sending entity. Second, the
   importance of authentication varies when different identities of a
   request are concerned. This specification defines new identity
   conditions that can accommodate the granularity of specific SIP
   identity header fields. The requirement for authentication depends
   on which field is to be matched.

   The identity condition for load filtering is specified by the
   <call-identity> element and its sub-element <sip>. The <sip> element
   itself contains sub-elements representing SIP sending and receiving
   identity header fields: <from>, <to>, <request-uri>, and
   <p-asserted-identity>. All those sub-elements are of an extended
   form of the [RFC4745] <identity> element. In addition to the sub-
   elements including <one>, <except>, and <many> in the <identity>
   element from [RFC4745], the extended form adds two new sub-elements,
   namely, <many-tel> and <except-tel>, which will be explained later in
   this section.

   The [RFC4745] <one> and <except> elements may contain an "id"
   attribute, which is the URI of a single entity to be included or
   excluded in the condition. When used in the <from>, <to>,
   <request-uri>, and <p-asserted-identity> elements, this "id" value is
   the URI contained in the corresponding SIP header field, i.e., From,
   To, Request-URI, and P-Asserted-Identity.

   When the <call-identity> element contains multiple <sip> sub-
   elements, the result is combined using logical OR. When the <from>,
   <to>, <request-uri>, and <p-asserted-identity> elements contain
   multiple <one>, <many>, or <many-tel> sub-elements, the result is
   also combined using logical OR. When the <many> sub-element further
   contains one or more <except> sub-elements, or when the <many-tel>
   sub-element further contains one or more <except-tel> sub-elements,
   the result of each <except> or <except-tel> sub-element is combined
   using a logical OR, similar to that of the [RFC4745] <identity>
   element. However, when the <sip> element contains multiple <from>,
   <to>, <request-uri>, and <p-asserted-identity> sub-elements, the

result is combined using logical AND. This allows the call identity
to be specified by multiple fields of a SIP request simultaneously,
e.g., both the From and the To header fields.

The following shows an example of the <call-identity> element, which
matches call requests whose To header field contains the SIP URI
"sip:alice@hotline.example.com" or the 'tel' URI
"tel:+1-212-555-1234".

```
          <call-identity>
             <sip>
                <to>
                    <one id="sip:alice@hotline.example.com"/>
                    <one id="tel:+1-212-555-1234"/>
                </to>
             </sip>
          </call-identity>
```

Before evaluating <call-identity> conditions, the subscriber shall
convert URIs received in SIP header fields in canonical form as per
[RFC3261], except that the "phone-context" parameter shall not be
removed, if present.

The [RFC4745] <many> and <except> elements may take a "domain"
attribute. The "domain" attribute specifies a domain name to be
matched by the domain part of the candidate identity. Thus, it
allows matching a large and possibly unknown number of entities
within a domain. The "domain" attribute works well for SIP URIs.

A URI identifying a SIP user, however, can also be a 'tel' URI.
Therefore, a similar way to match a group of 'tel' URIs is needed.
There are two forms of 'tel' URIs: for global numbers and local
numbers. According to [RFC3966], "All phone numbers MUST use the
global form unless they cannot be represented as such...Local numbers
MUST be tagged with a 'phone-context'". The global number 'tel' URIs
start with a "+". The "phone-context" parameter of local numbers may
be labeled as a global number or any number of its leading digits or
a domain name. Both forms of the 'tel' URI make the resulting URI
globally unique.

'tel' URIs of global numbers can be grouped by prefixes consisting of
any number of common leading digits. For example, a prefix formed by
a country code or both the country and area code identifies telephone
numbers within a country or an area. Since the length of the country
and area code for different regions are different, the length of the
number prefix also varies. This allows further flexibility such as

grouping the numbers into sub-areas within the same area code. 'tel'
URIs of local numbers can be grouped by the value of the
"phone-context" parameter.

The <many> and <except> sub-elements in the <identity> element of
[RFC4745] do not allow additional attributes to be added directly.
Redefining behavior of their existing "domain" attribute creates
backward-compatibility issues. Therefore, this specification defines
the <many-tel> and <except-tel> sub-elements that extend the
[RFC4745] <identity> element. Both of them have a "prefix" attribute
for grouping 'tel' URIs, similar to the "domain" attribute for
grouping SIP URIs in existing <many> and <except> sub-elements. For
global numbers, the "prefix" attribute value holds any number of
common leading digits, for example, "+1-212" for US phone numbers
within area code "212" or "+1-212-854" for the organization with US
area code "212" and local prefix "854". For local numbers, the
"prefix" attribute value contains the "phone-context" parameter
value. It should be noted that visual separators (such as the "-"
sign) in 'tel' URIs are not used for URI comparison as per [RFC3966].

The following example shows the use of the "prefix" attribute along
with the "domain" attribute. It matches those requests calling to
the number "+1-202-999-1234" but are not calling from a "+1-212"
prefix or a SIP From URI domain of "manhattan.example.com".

```
            <call-identity>
               <sip>
                   <from>
                       <many>
                           <except domain="manhattan.example.com"/>
                       </many>
                       <many-tel>
                           <except-tel prefix="+1-212"/>
                       </many-tel>
                   </from>
                   <to>
                       <one id="tel:+1-202-999-1234"/>
                   </to>
               </sip>
            </call-identity>
```

5.3.2. Method

The load created on a SIP server depends on the type of initial SIP
requests for dialogs or standalone transactions. The <method>
element specifies the SIP method to which the load-filtering action
applies. When this element is not included, the load-filtering
actions are applicable to all applicable initial requests. These

requests include INVITE, MESSAGE, REGISTER, SUBSCRIBE, OPTIONS, and
PUBLISH. Non-initial requests, such as ACK, BYE, and CANCEL MUST NOT
be subjected to load filtering. In addition, SUBSCRIBE requests are
not filtered if the event-type header field indicates the event
package defined in this specification.

The following example shows the use of the <method> element in the
case the filtering actions should be applied to INVITE requests.

        <method>INVITE</method>

5.3.3. Target SIP Entity

A SIP server that performs load-filtering may have multiple paths to
route call requests matching the same set of call identity elements.
In those situations, the SIP load-filtering server may desire to take
advantage of alternative paths and only apply load-filtering actions
to matching requests for the next-hop SIP entity that originated the
corresponding load-filtering policy. To achieve that, the SIP load-
filtering server needs to associate every load-filtering policy with
its originating SIP entity. The <target-sip-entity> element is
defined for that purpose, and it contains the URI of the entity that
initiated the load-filtering policy, which is generally the
corresponding notifier. A notifier MAY include this element as part
of the condition of its filtering policy being sent to the
subscriber, as below.

<target-sip-entity>sip:biloxi.example.com</target-sip-entity>

When a SIP load-filtering server receives a policy with a
<target-sip-entity> element, it SHOULD record it and take it into
consideration when making load-filtering decisions. If the load-
filtering server receives a load-filtering policy that does not
contain a <target-sip-entity> element, it MAY still record the URI of
the load-filtering policy's originator as the <target-sip-entity>
information and consider it when making load-filtering decisions.

   The following are two examples of using the <target-sip-entity>
   element.

   Use case I: The network has user A connected to SIP Proxy 1 (SP1),
   user B connected to SIP Proxy 3 (SP3), SP1 and SP3 connected via
   SIP Proxy 2 (SP2), and SP2 connected to an Application Server
   (AS). Under normal load conditions, a call from A to B is routed
   along the following path: A-SP1-SP2-AS-SP3-B. The AS provides a
   nonessential service and can be bypassed in case of overload. Now
   let's assume that AS is overloaded and sends to SP2 a load-
   filtering policy requesting that 50% of all INVITE requests be

dropped. SP2 can maintain AS as the <target-sip-entity> for that
policy so that it knows the 50% drop action is only applicable to
call requests that must go through AS, without affecting those
calls directly routed through SP3 to B.

Use case II: A translation service for toll-free numbers is
installed on two Application Servers, AS1 and AS2. User A is
connected to SP1 and calls 800-1234-4529, which is translated by
AS1 and AS2 into a regular E.164 number depending on, e.g., the
caller's location. SP1 forwards INVITE requests with Request-URI
= "800 number" to AS1 or AS2 based on a load-balancing strategy.
As calls to 800-1234-4529 create a pre-overload condition in AS1,
AS1 sends to SP1 a load-filtering policy requesting that 50% of
calls towards 800-1234-4529 be rejected. In this case, SP1 can
maintain AS1 as the <target-sip-entity> for the rule, and only
apply the load-filtering policy on incoming requests that are
intended to be sent to AS1. Those requests that are sent to AS2,
although matching the <call-identity> of the filter, will not be
affected.

5.3.4. Validity

A filtering policy is usually associated with a validity period
condition. This specification reuses the <validity> element of
[RFC4745], which specifies a period of validity time by pairs of
<from> and <until> sub-elements. When multiple time periods are
defined, the validity condition is evaluated to TRUE if the current
time falls into any of the specified time periods. That is, it
represents a logical OR operation across all validity time periods.

The following example shows a <validity> element specifying a valid
period from 12:00 to 15:00 US Eastern Standard Time on 2008-05-31.

```
<validity>
    <from>2008-05-31T12:00:00-05:00</from>
    <until>2008-05-31T15:00:00-05:00</until>
</validity>
```

5.4. Actions

The actions a load-filtering server takes on loads matching the load-
filtering conditions are defined by the <accept> element in the load-
filtering policy, which includes any one of the three sub-elements
<rate>, <percent>, and <win>. The <rate> element denotes an absolute
value of the maximum acceptable request rate in requests per second;
the <percent> element specifies the relative percentage of incoming
requests that should be accepted; the <win> element describes the
acceptable window size supplied by the receiver, which is applicable

   in window-based load-filtering. In static load-filtering policy
   configuration scenarios, using the <rate> sub-element is RECOMMENDED
   because it is hard to enforce the percentage rate or window-based
   load filtering when incoming load from upstream or reactions from
   downstream are uncertain. (See [SIP-OVERLOAD] and [RFC6357] for more
   details on rate-based, loss-based, and window-based load control.)

   In addition, the <accept> element takes an optional "alt-action"
   attribute that can be used to explicitly specify the desired action
   in case a request cannot be processed. The "alt-action" can take one
   of the following three values: "reject", "redirect", or "drop".

   o The "reject" action is the default value for "alt-action". It
      means that the load-filtering server will reject the request with
      a 503 "Service Unavailable" response message.

   o The "redirect" action means redirecting the request to another
      target. When it is used, an "alt-target" attribute MUST be
      defined. The "alt-target" specifies one URI or a list of URIs
      where the request should be redirected. The server sends out the
      redirect URIs in a 300-class response message.

   o The "drop" action means simply ignoring the request without doing
      anything, which can, in certain cases, help save processing
      capability during overload. For example, when SIP is running over
      a reliable transport such as TCP, the "drop" action does not send
      out the rejection response, neither does it close the transport
      connection. However, when running SIP over an unreliable
      transport such as UDP, using the "drop" action will create message
      retransmissions that further worsen the possible overload
      situation. Therefore, any "drop" action applied to an unreliable
      transport MUST be treated as if it were "reject".

   The above "alt-action" processing can also be illustrated through the
   following pseudocode.

```
        SWITCH "alt-action"
          "redirect": "redirect"
          "drop":
            IF unreliable-transport
              THEN treat as "reject"
            ELSE
              "drop"
          "reject": "reject"
          default: "reject"
        END
```

In the following <actions> element example, the server accepts
maximum of 100 call requests per second. The remaining calls are
redirected to an answering machine.

```
        <actions>
           <accept alt-action="redirect" alt-target=
                  "sip:answer-machine@example.com">
              <rate>100</rate>
           </accept>
        </actions>
```

6. XML Schema Definition for Load Control

   This section defines the XML schema for the load-control document.
   It extends the Common Policy schema in [RFC4745] in two ways.
   Firstly, it defines two mandatory attributes for the <ruleset>
   element: "version" and "state". The "version" attribute allows the
   recipient of the notification to properly order them. Versions start
   at zero and increase by one for each new document sent to a
   subscriber within the same subscription. Versions MUST be
   representable using a non-negative 32-bit integer. The "state"
   attribute indicates whether the document contains a full load-
   filtering policy update or only state delta as partial update.
   Secondly, it defines new members of the <conditions> and <actions>
   elements.

```
   <?xml version="1.0" encoding="UTF-8"?>
   <xs:schema targetNamespace="urn:ietf:params:xml:ns:load-control"
       xmlns:lc="urn:ietf:params:xml:ns:load-control"
       xmlns:cp="urn:ietf:params:xml:ns:common-policy"
       xmlns:xs="http://www.w3.org/2001/XMLSchema"
       elementFormDefault="qualified"
       attributedFormDefault="unqualified">

   <xs:import namespace="urn:ietf:params:xml:ns:common-policy"/>

   <!-- RULESET -->

   <xs:element name="ruleset">
     <xs:complexType>
       <xs:complexContent>
         <xs:restriction base="xs:anyType">
           <xs:sequence>
             <xs:element name="rule" type="cp:ruleType"
             minOccurs="0" maxOccurs="unbounded"/>
           </xs:sequence>
         </xs:restriction>
       </xs:complexContent>
```

```
        <xs:attribute name="version" type="xs:integer" use="required"/>
        <xs:attribute name="state" use="required">
          <xs:simpleType>
            <xs:restriction base="xs:string">
              <xs:enumeration value="full"/>
              <xs:enumeration value="partial"/>
            </xs:restriction>
          </xs:simpleType>
        </xs:attribute>
      </xs:complexType>
    </xs:element>

    <!-- CONDITIONS -->

    <!-- CALL IDENTITY -->
    <xs:element name="call-identity" type="lc:call-identity-type"/>

    <!-- CALL IDENTITY TYPE -->
    <xs:complexType name="call-identity-type">
      <xs:choice>
        <xs:element name="sip" type="lc:sip-id-type"/>
        <any namespace="##other" processContents="lax" minOccurs="0"
        maxOccurs="unbounded"/>
      </xs:choice>
      <anyAtrribute namespace="##other" processContents="lax"/>
    </xs:complexType>

    <!-- SIP ID TYPE -->
    <xs:complexType name="sip-id-type">
      <xs:sequence>
        <element name="from" type="lc:identityType" minOccurs="0"/>
        <element name="to" type="lc:identityType" minOccurs="0"/>
        <element name="request-uri" type="lc:identityType"
        minOccurs="0"/>
        <element name="p-asserted-identity" type="lc:identityType"
        minOccurs="0"/>
        <any namespace="##other" processContents="lax" minOccurs="0"
        maxOccurs="unbounded"/>
      </xs:sequence>
      <anyAtrribute namespace="##other" processContents="lax"/>
    </xs:complexType>

    <!-- IDENTITY TYPE -->
    <xs:complexType name="identityType">
      <xs:complexContent>
        <xs:restriction base="xs:anyType">
          <xs:choice minOccurs="1" maxOccurs="unbounded">
            <xs:element name="one" type="cp:oneType"/>
```

```
            <xs:element name="many" type="lc:manyType"/>
            <xs:element name="many-tel" type="lc:manyTelType"/>
            <xs:any namespace="##other" processContents="lax"/>
          </xs:choice>
        </xs:restriction>
      </xs:complexContent>
    </xs:complexType>

    <!-- MANY-TEL TYPE -->
    <xs:complexType name="manyTelType">
      <xs:complexContent>
        <xs:restriction base="xs:anyType">
          <xs:choice minOccurs="0" maxOccurs="unbounded">
            <xs:element name="except-tel" type="lc:exceptTelType"/>
            <xs:any namespace="##other"
            minOccurs="0" processContents="lax"/>
          </xs:choice>
          <xs:attribute name="prefix"
          use="optional" type="xs:string"/>
        </xs:restriction>
      </xs:complexContent>
    </xs:complexType>

    <!-- EXCEPT-TEL TYPE -->
    <xs:complexType name="exceptTelType">
      <xs:attribute name="prefix" type="xs:string" use="optional"/>
      <xs:attribute name="id" type="xs:anyURI" use="optional"/>
    </xs:complexType>

    <!-- METHOD -->
    <xs:element name="method" type="lc:method-type"/>

    <!-- METHOD TYPE -->
    <xs:simpleType name="method-type">
      <xs:restriction base="xs:string">
        <xs:enumeration value="INVITE"/>
        <xs:enumeration value="MESSAGE"/>
        <xs:enumeration value="REGISTER"/>
        <xs:enumeration value="SUBSCRIBE"/>
        <xs:enumeration value="OPTIONS"/>
        <xs:enumeration value="PUBLISH"/>
      </xs:restriction>
    </xs:simpleType>

    <!-- TARGET SIP ENTITY -->
    <xs:element name="target-sip-entity" type="xs:anyURI" minOccurs="0"/>

    <!-- ACTIONS -->
```

```
      <xs:element name="accept">
        <xs:choice>
          <element name="rate" type="xs:decimal" minOccurs="0"/>
          <element name="win" type="xs:integer" minOccurs="0"/>
          <element name="percent" type="xs:decimal" minOccurs="0"/>
          <any namespace="##other" processContents="lax" minOccurs="0"
          maxOccurs="unbounded"/>
        </xs:choice>
        <xs:attribute name="alt-action" type="xs:string" default="reject"/>
        <xs:attribute name="alt-target" type="lc:alt-target-type"
        use="optional"/>
        <anyAtrribute namespace="##other" processContents="lax"/>
      </xs:element>

      <!-- ALT TARGET TYPE -->
      <xs:simpleType name="alt-target-type">
        <xs:list itemType="xs:anyURI"/>
      </xs:simpleType>

      </xs:schema>
```

7. Security Considerations

   Two primary security considerations arise from this specification.
   One is the distribution mechanism for the load filtering policy that
   is based on the SIP event notification framework, and the other is
   the enforcement mechanism for the load-filtering policy.

   Security considerations for SIP event package mechanisms are covered
   in Section 6 of [RFC6665]. A particularly relevant security concern
   for this event package is that if the notifiers can be spoofed,
   attackers can send fake notifications asking subscribers to throttle
   all traffic, leading to denial-of-service (DoS) attacks. Therefore,
   this SIP load-filtering mechanism MUST be used in a Trust Domain
   (Section 3.4). But if a legitimate notifier in the Trust Domain is
   itself compromised, additional mechanisms will be needed to detect
   the attack.

   Security considerations for load-filtering policy enforcement depends
   very much on the contents of the policy. This specification defines
   a possible match of the following SIP header fields in a load-
   filtering policy: <from>, <to>, <request-uri>, and
   <p-asserted-identity>. The exact requirement to authenticate and
   authorize these fields is up to the service provider. In general, if
   the identity field represents the source of the request, it SHOULD be
   authenticated and authorized; if the identity field represents the
   destination of the request, the authentication and authorization is
   optional.

In addition, the "redirect" action (Section 5.4) could facilitate a
reflection denial-of-service attack. If a number of SIP proxy
servers in a Trust Domain are using UDP and configured to get their
policies from a central server. An attacker spoofs the central
server's address to send a number of NOTIFY bodies telling the proxy
servers to redirect all calls to victim@outside-of-trust-domain.com.
The proxy servers then redirect all calls to the victim, who then
becomes a victim of Denial of Service attack and becomes
inaccessible from the Internet. To address this type of threat,
this specification requires that a Trust Domain agrees on what types
of calls can be affected as well as on the destinations to which
calls may be redirected, as in Section 3.4.

8. IANA Considerations

   This specification registers a SIP event package, a new media type, a
   new XML namespace, and a new XML schema.

8.1. Load-Control Event Package Registration

   This section registers an event package based on the registration
   procedures defined in [RFC6665].

   Package name: load-control

   Type: package

   Published specification: This specification

   Person to contact: Charles Shen, charles@cs.columbia.edu

8.2. application/load-control+xml Media Type Registration

   This section registers a new media type based on the procedures
   defined in [RFC6838] and guidelines in [RFC3023].

   Type name: application

   Subtype name: load-control+xml

   Required parameters: none

   Optional parameters: Same as charset parameter of application/xml as
   specified in [RFC3023].

   Encoding considerations: Same as encoding considerations of
   application/xml as specified in [RFC3023].

   Security considerations: See Section 10 of [RFC3023] and Section 7 of
   this specification.

   Interoperability considerations: none

   Published specification: This specification

   Applications that use this media type: Applications that perform load
   control of SIP entities.

   Fragment identifier considerations: Same as fragment identifier
   considerations of application/xml as specified in [RFC3023].

   Additional Information:

      Deprecated alias names for this type: none

      Magic Number(s): none

      File Extension(s): .xml

      Macintosh file type code(s): "TEXT"

   Person and email address for further information: Charles Shen,
   charles@cs.columbia.edu

   Intended usage: COMMON

   Restrictions on usage: none

   Author: Charles Shen, Henning Schulzrinne, Arata Koike

   Change controller: IESG

   Provisional registration? (standards tree only): no

8.3. URN Sub-Namespace Registration

   This section registers a new XML namespace, as per the guidelines in
   [RFC3688]

   URI: The URI for this namespace is

      urn:ietf:params:xml:ns:load-control

   Registrant Contact: IETF SOC Working Group <sip-overload@ietf.org>,
   as designated by the IESG <iesg@ietf.org>

XML:

```
BEGIN
<?xml version="1.0"?>
<!DOCTYPE html PUBLIC "-//W3C//DTD XHTML Basic 1.0//EN"
             "http://www.w3.org/TR/xhtml-basic/xhtml-basic10.dtd">
<html xmlns="http://www.w3.org/1999/xhtml">
<head>
  <meta http-equiv="content-type"
    content="text/html;charset=iso-8859-1"/>
  <title>SIP Load-Control Namespace</title>
</head>
<body>
  <h1>Namespace for SIP Load Control</h1>
  <h2>urn:ietf:params:xml:ns:load-control</h2>
  <p>See <a href="http://www.rfc-editor.org/rfc/rfc7200.txt">
     RFC 7200</a>.</p>
</body>
</html>
END
```

8.4. Load-Control Schema Registration

URI: urn:ietf:params:xml:schema:load-control

Registrant Contact: IETF SOC working group, Charles Shen
(charles@cs.columbia.edu).

XML: the XML schema contained in Section 6 has been registered.

Its first line is

<?xml version="1.0" encoding="UTF-8"?>

and its last line is

</xs:schema>

9. Acknowledgements

10. References

10.1. Normative References

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10.2. Informative References

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   [RFC4412] Schulzrinne, H. and J. Polk, "Communications Resource
              Priority for the Session Initiation Protocol (SIP)", RFC
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   [RFC4825] Rosenberg, J., "The Extensible Markup Language (XML)
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   [RFC6357] Hilt, V., Noel, E., Shen, C., and A. Abdelal, "Design
              Considerations for Session Initiation Protocol (SIP)
              Overload Control", RFC 6357, August 2011.

   [SIP-OVERLOAD]
              Gurbani, V., Ed., Hilt, V., and H. Schulzrinne, "Session
              Initiation Protocol (SIP) Overload Control", Work in
              Progress, March 2014.

Appendix A. Definitions

   This specification reuses the definitions for "Event Package",
   "Notification", "Notifier", "Subscriber", and "Subscription" as in
   [RFC6665]. The following additional definitions are also used.

   Load Filtering: A load-control mechanism that applies specific
      actions to selected loads (e.g., SIP requests) matching specific
      conditions.

   Load-Filtering Policy: A set of zero or more load-filtering rules,
      also known as load-filtering rule set.

   Load-Filtering Rule: Conditions and actions to be applied for load
      filtering.

   Load-Filtering Condition: Elements that describe how to select loads
      to apply load-filtering actions. This specification defines the
      <call-identity>, <method>, <target-sip-identity>, and <validity>
      condition elements (Section 5.3).

   Load-Filtering Action: An operation to be taken by a load-filtering
      server on loads that match the load-filtering conditions. This
      specification allows actions such as accept, reject, and redirect
      of loads (Section 5.4).

   Load-Filtering Server: A server that performs load filtering. In
      the context of this specification, the load-filtering server is
      the subscriber, which receives load-filtering policies from the
      notifier and enforces those policies during load filtering.

   Load-Control Document: An XML document that describes the load-
      filtering policies (Section 5). It inherits and enhances the
      common policy document defined in [RFC4745].

Appendix B. Design Requirements

   The SIP load-filtering mechanism needs to satisfy the following
   requirements:

   o For simplicity, the solution should focus on a method for
      controlling SIP load, rather than a generic application-layer
      mechanism.

   o The load-filtering policy needs to be distributed efficiently to
      possibly a large subset of all SIP elements.

   o The solution should reuse existing SIP protocol mechanisms to
      reduce implementation and deployment complexity.

   o For predictable overload situations, such as holidays and mass
      calling events, the load-filtering policy should specify during
      what time it is to be applied, so that the information can be
      distributed ahead of time.

   o For destination-specific overload situations, the load-filtering
      policy should be able to describe the destination domain or the
      callee.

   o To address accidental and intentional high-volume call generators,
      the load-filtering policy should be able to specify the caller.

   o Caller and callee need to be specified as both SIP URIs and 'tel'
      URIs [RFC3966] in load-filtering policies.

   o It should be possible to specify particular information in the SIP
      headers (e.g., prefixes in telephone numbers) that allow load
      filtering over limited regionally focused overloads.

   o The solution should draw upon experiences from related PSTN
      mechanisms [Q.1248.2] [E.412] [E.300SerSup3] where applicable.

   o The solution should be extensible to meet future needs.

Appendix C. Discussion of How This Specification Meets the Requirements
             of RFC 5390

   This section evaluates whether the load-control event package
   mechanism defined in this specification satisfies various SIP
   overload control requirements set forth by [RFC5390]. As mentioned
   in Section 1, this specification complements other efforts in the
   overall SIP load-control solution space. Therefore, not all RFC 5390
   requirements are found applicable to this specification. This
   specification categorizes the assessment results into Yes (the
   requirement is met), P/A (Partially Applicable), No (must be used in
   conjunction with another mechanism to meet the requirement), and N/A
   (Not Applicable).

      REQ 1: The overload mechanism shall strive to maintain the overall
      useful throughput (taking into consideration the quality-of-
      service needs of the using applications) of a SIP server at
      reasonable levels, even when the incoming load on the network is
      far in excess of its capacity. The overall throughput under load
      is the ultimate measure of the value of an overload control
      mechanism.

   P/A. The goal of load filtering is to prevent overload or maintain
   overall goodput during the time of overload, but it is dependent on
   the predictions of the load and the computations as well as
   distribution of the filtering policies. If the load predictions or
   filtering policy computations are incorrect, or the filtering policy
   is not properly distributed, the mechanism will be less effective.
   On the other hand, if the load can be accurately predicted and
   filtering policies be computed and distributed appropriately, this
   requirement can be met.

      REQ 2: When a single network element fails, goes into overload, or
      suffers from reduced processing capacity, the mechanism should
      strive to limit the impact of this on other elements in the
      network. This helps to prevent a small-scale failure from
      becoming a widespread outage.

   N/A if load-filtering policies are installed in advance and do not
   change during the potential overload period, P/A if load-filtering
   policies are dynamically adjusted. The algorithm to dynamically
   compute load-filtering policies is outside the scope of this
   specification, while the distribution of the updated filtering
   policies uses the event package mechanism of this specification.

      REQ 3: The mechanism should seek to minimize the amount of
      configuration required in order to work. For example, it is
      better to avoid needing to configure a server with its SIP message
      throughput, as these kinds of quantities are hard to determine.

   No. This mechanism is entirely dependent on advance configuration,
   based on advance knowledge. In order to satisfy REQ 3, it should be
   used in conjunction with other mechanisms that are not based on
   advance configuration.

      REQ 4: The mechanism must be capable of dealing with elements that
      do not support it, so that a network can consist of a mix of
      elements that do and don't support it. In other words, the
      mechanism should not work only in environments where all elements
      support it. It is reasonable to assume that it works better in
      such environments, of course. Ideally, there should be
      incremental improvements in overall network throughput as
      increasing numbers of elements in the network support the
      mechanism.

   No. This mechanism is entirely dependent on the participation of all
   possible neighbors. In order to satisfy REQ 4, it should be used in
   conjunction with other mechanisms, some of which are described in
   Section 3.4.

REQ 5: The mechanism should not assume that it will only be
deployed in environments with completely trusted elements. It
should seek to operate as effectively as possible in environments
where other elements are malicious; this includes preventing
malicious elements from obtaining more than a fair share of
service.

No. This mechanism is entirely dependent on the non-malicious
participation of all possible neighbors. In order to satisfy REQ 5,
it should be used in conjunction with other mechanisms, some of which
are described in Section 3.4.

REQ 6: When overload is signaled by means of a specific message,
the message must clearly indicate that it is being sent because of
overload, as opposed to other, non overload-based failure
conditions. This requirement is meant to avoid some of the
problems that have arisen from the reuse of the 503 response code
for multiple purposes. Of course, overload is also signaled by
lack of response to requests. This requirement applies only to
explicit overload signals.

N/A. This mechanism signals anticipated overload, not actual
overload. However, the signals in this mechanism are not used for
any other purpose.

REQ 7: The mechanism shall provide a way for an element to
throttle the amount of traffic it receives from an upstream
element. This throttling shall be graded so that it is not all-
or-nothing as with the current 503 mechanism. This recognizes the
fact that "overload" is not a binary state and that there are
degrees of overload.

Yes. This event package allows rate-/loss-/window-based overload
control options as discussed in Section 5.4.

REQ 8: The mechanism shall ensure that, when a request was not
processed successfully due to overload (or failure) of a
downstream element, the request will not be retried on another
element that is also overloaded or whose status is unknown. This
requirement derives from REQ 1.

N/A to the load-control event package mechanism itself.

REQ 9: That a request has been rejected from an overloaded element
shall not unduly restrict the ability of that request to be
submitted to and processed by an element that is not overloaded.
This requirement derives from REQ 1.

Yes. For example, load-filtering policy (Section 3.1) can include
alternative forwarding destinations for rejected requests.

    REQ 10: The mechanism should support servers that receive requests
    from a large number of different upstream elements, where the set
    of upstream elements is not enumerable.

No. Because this mechanism requires advance configuration of
specifically identified neighbors, it does not support environments
where the number and identity of the upstream neighbors are not known
in advance. In order to satisfy REQ 10, it should be used in
conjunction with other mechanisms.

    REQ 11: The mechanism should support servers that receive requests
    from a finite set of upstream elements, where the set of upstream
    elements is enumerable.

Yes. See also answer to REQ 10.

    REQ 12: The mechanism should work between servers in different
    domains.

Yes. The load-control event package mechanism is not limited by
domain boundaries. However, it is likely more applicable in intra-
domain scenarios than in inter-domain scenarios due to security and
other concerns (see also Section 3.4).

    REQ 13: The mechanism must not dictate a specific algorithm for
    prioritizing the processing of work within a proxy during times of
    overload. It must permit a proxy to prioritize requests based on
    any local policy, so that certain ones (such as a call for
    emergency services or a call with a specific value of the
    Resource-Priority header field [RFC4412]) are given preferential
    treatment, such as not being dropped, being given additional
    retransmission, or being processed ahead of others.

P/A. This mechanism does not specifically address the prioritizing
of work during times of overload. But it does not preclude any
particular local policy.

    REQ 14: The mechanism should provide unambiguous directions to
    clients on when they should retry a request and when they should
    not. This especially applies to TCP connection establishment and
    SIP registrations, in order to mitigate against avalanche restart.

N/A to the load-control event package mechanism itself.

      REQ 15: In cases where a network element fails, is so overloaded
      that it cannot process messages, or cannot communicate due to a
      network failure or network partition, it will not be able to
      provide explicit indications of the nature of the failure or its
      levels of congestion. The mechanism must properly function in
      these cases.

   P/A. Because the load-filtering policies are provisioned in advance,
   they are not affected by the overload or failure of other network
   elements. On the other hand, they may not, in those cases, be able
   to protect the overloaded network elements (see REQ 1).

      REQ 16: The mechanism should attempt to minimize the overhead of
      the overload control messaging.

   Yes. The standardized SIP event package mechanism [RFC6665] is used.

      REQ 17: The overload mechanism must not provide an avenue for
      malicious attack, including DoS and DDoS attacks.

   P/A. This mechanism does provide a potential avenue for malicious
   attacks. Therefore, the security mechanisms for SIP event packages,
   in general, [RFC6665] and Section 7 of this specification should be
   used.

      REQ 18: The overload mechanism should be unambiguous about whether
      a load indication applies to a specific IP address, host, or URI,
      so that an upstream element can determine the load of the entity
      to which a request is to be sent.

   Yes. The identity of load indication is covered in the load-
   filtering policy format definition in Section 3.1.

      REQ 19: The specification for the overload mechanism should give
      guidance on which message types might be desirable to process over
      others during times of overload, based on SIP-specific
      considerations. For example, it may be more beneficial to process
      a SUBSCRIBE refresh with Expires of zero than a SUBSCRIBE refresh
      with a non-zero expiration (since the former reduces the overall
      amount of load on the element), or to process re-INVITEs over new
      INVITEs.

   N/A to the load-control event package mechanism itself.

      REQ 20: In a mixed environment of elements that do and do not
      implement the overload mechanism, no disproportionate benefit
      shall accrue to the users or operators of the elements that do not
      implement the mechanism.

   No. This mechanism is entirely dependent on the participation of all
   possible neighbors. In order to satisfy REQ 20, it should be used in
   conjunction with other mechanisms, some of which are described in
   Section 3.4.

      REQ 21: The overload mechanism should ensure that the system
      remains stable. When the offered load drops from above the
      overall capacity of the network to below the overall capacity, the
      throughput should stabilize and become equal to the offered load.

   N/A to the load-control event package mechanism itself.

      REQ 22: It must be possible to disable the reporting of load
      information towards upstream targets based on the identity of
      those targets. This allows a domain administrator who considers
      the load of their elements to be sensitive information, to
      restrict access to that information. Of course, in such cases,
      there is no expectation that the overload mechanism itself will
      help prevent overload from that upstream target.

   N/A to the load-control event package mechanism itself.

      REQ 23: It must be possible for the overload mechanism to work in
      cases where there is a load balancer in front of a farm of
      proxies.

   Yes. The load-control event package mechanism does not preclude its
   use in a scenario with server farms.

Appendix D. Complete Examples

D.1. Load-Control Document Examples

   This section presents two complete examples of load-control documents
   valid with respect to the XML schema defined in Section 6.

   The first example assumes that a set of hotlines are set up at
   "sip:alice@hotline.example.com" and "tel:+1-212-555-1234". The
   hotlines are activated from 12:00 to 15:00 US Eastern Standard Time
   on 2008-05-31. The goal is to limit the incoming calls to the
   hotlines to 100 requests per second. Calls that exceed the rate
   limit are explicitly rejected.

```
<?xml version="1.0" encoding="UTF-8"?>
<ruleset xmlns="urn:ietf:params:xml:ns:common-policy"
         xmlns:lc="urn:ietf:params:xml:ns:load-control"
         version="0" state="full">

    <rule id="f3g44k1">
        <conditions>
            <lc:call-identity>
                <lc:sip>
                    <lc:to>
                        <one id="sip:alice@hotline.example.com"/>
                        <one id="tel:+1-212-555-1234"/>
                    </lc:to>
                </lc:sip>
            </lc:call-identity>
            <method>INVITE</method>
            <validity>
                <from>2008-05-31T12:00:00-05:00</from>
                <until>2008-05-31T15:00:00-05:00</until>
            </validity>
        </conditions>
        <actions>
            <lc:accept alt-action="reject">
                <lc:rate>100</lc:rate>
            </lc:accept>
        </actions>

    </rule>
</ruleset>
```

The second example optimizes the usage of server resources during the
three-day period following a hurricane. Incoming calls to the domain
"sandy.example.com" or to call destinations with prefix "+1-212" will
be limited to a rate of 100 requests per second, except for those
calls originating from a particular rescue team domain
"rescue.example.com". Outgoing calls from the hurricane domain or
calls within the local domain are never limited. All calls that are
throttled due to the rate limit will be forwarded to an answering
machine with updated hurricane rescue information.

```
<?xml version="1.0" encoding="UTF-8"?>
<ruleset xmlns="urn:ietf:params:xml:ns:common-policy"
    xmlns:lc="urn:ietf:params:xml:ns:load-control"
    version="1" state="full">

    <rule id="f3g44k2">
        <conditions>
            <lc:call-identity>
                <lc:sip>
                    <lc:to>
                        <many domain="sandy.example.com"/>
                        <many-tel prefix="+1-212"/>
                    </lc:to>
                    <lc:from>
                        <many>
                            <except domain="sandy.example.com"/>
                            <except domain="rescue.example.com"/>
                        </many>
                    </lc:from>
                </lc:sip>
            </lc:call-identity>
            <method>INVITE</method>
            <validity>
                <from>2012-10-25T09:00:00+01:00</from>
                <until>2012-10-28T09:00:00+01:00</until>
            </validity>
        </conditions>
        <actions>
            <lc:accept alt-action="redirect" alt-target=
                    "sip:sandy@update.example.com">
                <lc:rate>100</lc:rate>
            </lc:accept>
        </actions>

    </rule>
</ruleset>
```

Sometimes it may occur that multiple rules in a ruleset define
actions that match the same methods, call identity and validity. In
those cases, the "first-match-wins" principle is used. For example,
in the following ruleset, the first rule requires all calls from the
"example.com" domain to be rejected. Even though the rule following
that one specifies that calls from "sip:alice@example.com" be
redirected to a specific target "sip:eve@example.com", the calls from
"sip:alice@example.com" will still be rejected because they have
already been matched by the earlier rule.

```
<?xml version="1.0" encoding="UTF-8"?>
<ruleset xmlns="urn:ietf:params:xml:ns:common-policy"
    xmlns:lc="urn:ietf:params:xml:ns:load-control"
    version="1" state="full">

    <rule id="f3g44k3">
        <conditions>
            <lc:call-identity>
                <lc:sip>
                    <lc:from>
                        <many domain="example.com"/>
                    </lc:from>
                </lc:sip>
            </lc:call-identity>
            <method>INVITE</method>
            <validity>
                <from>2013-7-2T09:00:00+01:00</from>
                <until>2013-7-3T09:00:00+01:00</until>
            </validity>
        </conditions>
        <actions>
            <lc:accept alt-action="reject">
                <lc:rate>0</lc:rate>
            </lc:accept>
        </actions>
    </rule>

    <rule id="f3g44k4">
        <conditions>
            <lc:call-identity>
                <lc:sip>
                    <lc:from>
                        <one id="sip:alice@example.com"/>
                    </lc:from>
                </lc:sip>
            </lc:call-identity>
            <method>INVITE</method>
            <validity>
                <from>2013-7-2T09:00:00+01:00</from>
                <until>2013-7-3T09:00:00+01:00</until>
            </validity>
        </conditions>
        <actions>
            <lc:accept alt-action="redirect" alt-target=
                    "sip:eve@example.com">
                <lc:rate>0</lc:rate>
            </lc:accept>
        </actions>
```

```
        </rule>
     </ruleset>
```

D.2. Message Flow Examples

   This section presents an example message flow of using the load-
   control event package mechanism defined in this specification.

```
        atlanta              biloxi
           | F1 SUBSCRIBE       |
           |------------------->|
           | F2 200 OK          |
           |<------------------ |
           | F3 NOTIFY          |
           |<------------------ |
           | F4 200 OK          |
           |------------------->|
```

        F1 SUBSCRIBE atlanta.example.com -> biloxi.example.com

           SUBSCRIBE sip:biloxi.example.com SIP/2.0
           Via: SIP/2.0/TCP atlanta.example.com;branch=z9hG4bKy7cjbu3
           From: sip:atlanta.example.com;tag=162ab5
           To: sip:biloxi.example.com
           Call-ID: 2xTb9vxSit55XU7p8@atlanta.example.com
           CSeq: 2012 SUBSCRIBE
           Contact: sip:atlanta.example.com
           Event: load-control
           Max-Forwards: 70
           Accept: application/load-control+xml
           Expires: 3600
           Content-Length: 0

        F2 200 OK   biloxi.example.com -> atlanta.example.com

           SIP/2.0 200 OK
           Via: SIP/2.0/TCP biloxi.example.com;branch=z9hG4bKy7cjbu3
             ;received=192.0.2.1
           To: <sip:biloxi.example.com>;tag=331dc8
           From: <sip:atlanta.example.com>;tag=162ab5
           Call-ID: 2xTb9vxSit55XU7p8@atlanta.example.com
           CSeq: 2012 SUBSCRIBE
           Expires: 3600
           Contact: sip:biloxi.example.com
           Content-Length: 0
```

```
   F3 NOTIFY  biloxi.example.com -> atlanta.example.com

      NOTIFY sip:atlanta.example.com SIP/2.0
      Via: SIP/2.0/TCP biloxi.example.com;branch=z9hG4bKy71g2ks
      From: <sip:biloxi.example.com>;tag=331dc8
      To: <sip:atlanta.example.com>;tag=162ab5
      Call-ID: 2xTb9vxSit55XU7p8@atlanta.example.com
      Event: load-control
      Subscription-State: active;expires=3599
      Max-Forwards: 70
      CSeq: 1775 NOTIFY
      Contact: sip:biloxi.example.com
      Content-Type: application/load-control+xml
      Content-Length: ...

      [Load-Control Document]

   F4 200 OK atlanta.example.com -> biloxi.example.com

      SIP/2.0 200 OK
      Via: SIP/2.0/TCP atlanta.example.com;branch=z9hG4bKy71g2ks
        ;received=192.0.2.2
      From: <sip:biloxi.example.com>;tag=331dc8
      To: <sip:atlanta.example.com>;tag=162ab5
      Call-ID: 2xTb9vxSit55XU7p8@atlanta.example.com
      CSeq: 1775 NOTIFY
      Content-Length: 0
```

Appendix E. Related Work

E.1. Relationship to Load Filtering in PSTN

   It is known that an existing PSTN network also uses a load-filtering
   mechanism to prevent overload and the filtering policy configuration
   is done manually except in specific cases when the Intelligent
   Network architecture is used [Q.1248.2][E.412]. This specification
   defines a load-filtering mechanism based on the SIP event
   notification framework that allows automated filtering policy
   distribution in suitable environments.

   PSTN overload control uses messages that specify an outgoing control
   list, call gap duration, and control duration [Q.1248.2][E.412].
   These items correspond roughly to the identity, action, and time
   fields of the SIP load-filtering policy defined in this
   specification. However, the load-filtering policy defined in this
   specification is much more generic and flexible as opposed to its
   PSTN counterpart.

   Firstly, PSTN load filtering only applies to telephone numbers. The
   identity element of SIP load-filtering policy allows both SIP URI and
   telephone numbers (through 'tel' URI) to be specified. These
   identities can be arbitrarily grouped by SIP domains or any number of
   leading prefixes of the telephone numbers.

   Secondly, the PSTN load-filtering action is usually limited to call
   gapping. The action field in SIP load-filtering policy allows more
   flexible possibilities such as rate throttle and others.

   Thirdly, the duration field in PSTN load filtering specifies a value
   in seconds for the load-filtering duration only, and the allowed
   values are mapped into a value set. The time field in SIP load-
   filtering policy may specify not only a duration, but also a future
   activation time that could be especially useful for automating load
   filtering for predictable overloads.

   PSTN load filtering can be performed in both edge switches and
   transit switches; the SIP load filtering can also be applied in both
   edge proxy servers and core proxy servers, and even in capable user
   agents.

   PSTN load filtering also has special accommodation for High
   Probability of Completion (HPC) calls, which would be similar to
   calls designated by the SIP Resource Priority Headers [RFC4412]. The
   SIP load-filtering mechanism also allows prioritizing the treatment
   of these calls by specifying favorable actions for them.

   PSTN load filtering also provides an administrative option for
   routing failed call attempts to either a reorder tone [E.300SerSup3]
   indicating overload conditions or a special recorded announcement. A
   similar capability can be provided in the SIP load-filtering
   mechanism by specifying appropriate "alt-action" attribute in the SIP
   load-filtering action field.

E.2. Relationship with Other IETF SIP Overload Control Efforts

   The load-filtering policies in this specification consist of
   identity, action, and time. The identity can range from a single
   specific user to an arbitrary user aggregate, domains, or areas. The
   user can be identified by either the source or the destination. When
   the user is identified by the source and a favorable action is
   specified, the result is, to some extent, similar to identifying a
   priority user based on authorized Resource Priority Headers [RFC4412]
   in the requests. Specifying a source user identity with an
   unfavorable action would cause an effect to some extent similar to an
   inverse SIP resource priority mechanism.

The load-filtering policy defined in this specification is generic
and expected to be applicable not only to the load-filtering
mechanism but also to the feedback overload control mechanism in
[SIP-OVERLOAD]. In particular, both mechanisms could use specific or
wildcard identities for load control and could share well-known load-
control actions. The time duration field in the load-filtering
policy could also be used in both mechanisms. As mentioned in
Section 1, the load-filtering policy distribution mechanism and the
feedback overload control mechanism address complementary areas in
the overload control problem space. Load filtering is more proactive
and focuses on distributing filtering policies towards the source of
the traffic; the hop-by-hop feedback-based approach is reactive and
reduces traffic already accepted by the network. Therefore, they
could also make different use of the generic load-filtering policy
components. For example, the load-filtering mechanism may use the
time field in the filtering policy to specify not only a control
duration but also a future activation time to accommodate a
predicable overload such as the one caused by Mother's Day greetings
or a viewer-voting program; the feedback-based control might not need
to use the time field or might use the time field to specify an
immediate load-control duration.

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